# Nick Weaver Fall 2018

### CS 161 Computer Security

## Midterm 1

Print your name:	(last)	(first)	
· ·	` '	dent Conduct and acknowledge the	ut academic misconduct
Sign your name:			
Print your class accou	nt login: cs161	and SID:	
Name of the person sitting to your left: —		Name of the person sitting to your right: —	
v	lectronic devices are not p	may not consult other notes, textlermitted. We use Gradescope for	
		up to the front of the exam room ing assumptions to the central doc	
	There are 11 questions, or old spending too long on a	of varying credit (134 points total) ny one question.	. The questions are of
Some of the test may in the footnotes.	nclude interesting technica	l asides as footnotes. You are not	responsible for reading
	Do not turn this page until	l your instructor tells you to do so	

Problem Ans		Cryptography True/False the following cryptography questions true or fa	alse.	(18 points)
(a)	(a) Let $E_k$ be a secure block cipher. True or False: It is impossible to find two messages $m$ and such that $m \neq m'$ and $E_k(m) = E_k(m')$ , even if the attacker knows $k$ .			
	0	True	0	False
(b)		$E_k$ be a secure block cipher. True or False $k$ ) and $(m',k')$ such that $m \neq m', k \neq k'$ and		
	0	True	0	False
(c)	Let MAC <sub>k</sub> be a secure MAC. TRUE or FALSE: It is computationally difficult to find messages $m$ and $m'$ such that $m \neq m'$ and MAC <sub>k</sub> $(m) = \text{MAC}_k(m')$ , even if the attacker knows $k$ .			
	0	True	0	FALSE
(d)		H be a cryptographic hash function. True sage $M$ .	or F	ALSE: $H(M)$ provides confidentiality for the
	0	True	0	FALSE
(e)	HM	IAC-DRBG does not have rollback resistance.		
	0	True	0	FALSE
(f)	Diff	Diffie/Hellman is secure in the presence of an active adversary.		
	0	True	0	FALSE
(g)	(g) Properly constructed RSA Signatures provide both integrity and authenticity.			tegrity and authenticity.
	0	True	0	FALSE
(h)	El (	Gamal encryption provides confidentiality but	it de	pes not provide integrity or authentication.
	0	True	0	FALSE
(i)		examining a certificate we need to consider how e's contents and signatures.	v we	obtained the certificate as well as the certifi-
	0	True	0	False

Proble	m 2	Potpourri		(18 points)
(a)	from	ead of storing user input on the stack, you de code, static, heap, and stack) for storing user memory address) of the section. Name or	er in	put. You also put a 64-bit canary at the top
(b)	Nam	e one memory-safety issue that the scheme fi	rom	part (a) fails to prevent.
(c)		E or FALSE: In a threat detection systems, fals always harmless.	e ne	gatives can be catastrophic, but false positives
	0	$\Gamma_{ m RUE}$	0	FALSE
(d)	Whic		o pr	otect a password database? (Select all that
		Salting Passwords		Using a Fast Hashing Function
		Encrypting Passwords		Using a Slow Hashing Function
(e)	(e) A heap overflow or use-after-free vulnerability can allow the attacker to overwrite the of an object (that is, the pointer at the start of a C++ object that points to the actu the function, basically a pointer to an array of function pointers). Can this bypass without additional information?		object that points to the actual methods for	
	0		0	No
(f)	At w	hat rank did Grace Hopper retire?		
( )	_	Lieutenant Colonel	0	Captain
	0	Rear Admiral	0	Brigadier General
sec		et remote server. When she needs to submit he	er ho	t she stores with her homework answers in a omework, she uses the MAC to check that her the key needed to generate the MAC. Which
	0	Integrity and Confidentiality	0	Authentication and Confidentiality
	0	Integrity and Authentication	0	Only Integrity

(h)	) Which of the following attacks can be used against a crypto system? (Select all that apply.)			
		Side-Channel		Chosen-ciphertext
		Rolling-regression		Rubber-Hose Cryptanalysis
		Chosen-plaintext		
(i)	"Cr	rypto" means:		
		Cryptography		Kryptonite
		Cryptocurrency		CryptoKitties
(j)	$Th\epsilon$	e Magic Word is:		
		Adava Kedavra		Stupify
		Windgardium Leviosa		Crucio

#### Problem 3 Security Principles

(12 points)

Write the best match for which security principle each situation.

Four CS 161 students, Chiyo, Habiba, Mr. Anderson, and Not Outis, decided that after learning about security principles and buffer overflows, they could implement their own distributed database (a database across multiple machines) with a focus on security!

- (a) Mr. Anderson suggests code their database in a higher-level programming language since they could avoid common security problems later on. Which security principle did he to use here?
- (b) Let's say they start coding their database and realized that a malicious user on one machine could corrupt their database. As a result, Habiba wants permission from at least 50% database users before a machine can be taken down. Which security principle is she using here?
- (c) The database the students built was password-protected for modification and they use a snippet (like the following) everywhere to check passwords:

String password = getPassword("user");

if (!password.equals(enteredPassword)) error();

Not Outis eventually forgets to put this snippet to check the passwords. What security principle does this violate?

- (d) To encrypt the data, Not Outis decides to take each piece of data and rotate the bytes in it by a fixed amount. It figured that since their database was closed source, no one would figure out how they were encrypting things. What security principle does this violate?
- (e) Mr. Anderson decides that new users should automatically get privileged access in order to set up their account to access whatever items they needed. After 1 hour, they would be dropped back to regular permissions, an administrator would be notified of changes, and they could revert changes if necessary. What security principle says this is not a good idea?
- (f) After fixing all previous problems, Chiyo decides to refactor their encryption code into its own module since a lot of it was spread across multiple modules. She also put all non-encryption code in a sandbox so that no vulnerabilities in those modules could effect the overall security of the database. What security principle is she trying to follow? What is she trying to minimize the size of?

The code below runs on a 32-bit Intel architecture. No defenses against buffer overflows are enabled. The code was not compiled to produce a position independent executable. No optimizations are enabled, and the compiler does not insert padding or reorder stack variables, which means buffer is at a lower address than fp.

```
int run_command(char *cmd) {
2
       return system (cmd);
3
4
  int print_hello(char *msg) {
       printf("Hello %s!\n", msg);
5
6
       return 0:
7
8
  int main() {
9
       int (*fp)(char *) = &print_hello;
10
       char buffer [8];
       gets (buffer);
11
12
       fp(buffer);
13
```

Note that the syntax int (\*fp)(char \*) indicates that fp is a pointer to a function which takes in a char \* and returns an int.

- (a) What line contains a memory vulnerability? What is this vulnerability called?
- (b) At line 12, we have that %ebp = Oxbfdead20 and &print\_hello = Ox08cafe13. Fill in the Python egg below to give an input which will overwrite the return address of main, causing the execution of the shellcode after the program returns from main.

print 'A' \* \_\_\_\_ + '\_\_\_\_\_' + 'AAAA' + '\_\_\_\_\_' ' + SHELLCODE

working? (Select all that apply.)

□ ASLR (same as part 5 on the project)

□ W^X

□ Using a memory-safe language instead of C

(c) Which of the following would sometimes or always prevent the code that you gave in part (b) from

(d) "I know," says Louis Reasoner, "let's add stack canaries to make this impossible to exploit!" Obviously this doesn't work. Fill in the Python egg below to give an input which will cause the execution of run\_command("/bin/sh"). At line 12, we have that %ebp = 0xbfdead20 and &run\_command = 0x08c0de42. Hint: Note that gets can read in a NUL byte (\x00), even in the middle of its input.

print '\_\_\_\_\_

- (e) Which of the following would sometimes or always prevent the code that you gave in part (d) from working? (Select all that apply.)
  - $\square$  ASLR (same as part 5 on the project)  $\square$  Selfrando
  - $\hfill\Box$  W^X  $\hfill\Box$  Using a memory-safe language instead of C

Ben Bitdiddle did not do a good job at coming up with a set of preconditions for some functions. For each code block, explain why with a short example the given preconditions are **not** sufficient to ensure memory safety by giving a small example.

```
(a)
       array_of_strings != NULL
 1
 2
       n \ll size(array_of_strings)
 3
       max_size > 0
 4
       for all i . 0 \ll i \ll n \implies
 5
            array\_of\_strings[i] != NULL and is a NUL-terminated string */
 6
   char *
 7
   concat_all(char *array_of_strings[], size_t n, size_t max_size) {
       char *concat = calloc(max_size, sizeof(char));
 9
       if (!concat) return NULL;
10
       size_t space_used = 0;
       for (size_t i = 0; i < n; i++) {
11
12
           char *s = array_of_strings[i];
13
            size_t len = strlen(s);
14
            strncpy(concat + space_used, s, max_size - space_used - 1);
15
            space_used += len;
16
17
       return concat;
18
```

Explanation:

```
(b)_{I}
        arr != NULL
 1
 2
        n \ll size(arr)
 3
        for \ all \ i \ . \ 0 <= i < n \Longrightarrow 0 <= arr[i] < n */
   int solve_interview_question(int *arr, size_t n) {
 4
 5
        for (size_t i = 0; i < n; i++)
 6
            arr[arr[i]] *= -1;
 7
        for (size_t i = 0; i < n; i++)
 8
            if (arr[i] < 0)
 9
                 return i;
10
        return 0;
11
```

Explanation:

(9 points)

Consider the code below.

```
void launch_nuclear_missiles() {
      puts("Launching the nukes...");
3
       /* code to launch nuclear missiles here */
4
       exit(1);
5
6
  #define MAX_INPUT 8
  int main() {
      char *correct_password = malloc(MAX_INPUT * sizeof(char));
10
      strcpy(correct_password, "S3creT\n");
11
      while (!feof(stdin)) {
12
           char *user_password = malloc(MAX_INPUT * sizeof(char));
           fgets (user_password, MAX_INPUT, stdin);
13
           if (strcmp(user_password, correct_password) == 0)
14
15
               launch_nuclear_missiles();
16
           free (user_password);
17
           free (correct_password);
18
           puts("Wrong password, try again!");
19
      }
20
```

All compiler optimizations are disabled, and both the source and binary are not available to David Lightman, who's trying to log in to play a game. Consider the following (buggy) interaction:

- 1. David inputs "Hello" followed by a newline.
- 2. The program outputs "Wrong password, try again!".
- 3. David inputs "Joshua" followed by a newline.
- 4. The program outputs "Launching the nukes...", and then the nukes are launched.<sup>1</sup>
- (a) Which memory safety vulnerability is present in this code?
- (b) Explain why this issue leads to the behavior David observes.

(c) How could you fix this issue in the code?

<sup>&</sup>lt;sup>1</sup>This immediately vaporizing millions of humans and wildlife on impact, beginning World War III and eventually wiping out most of the world due to an extended nuclear winter. This is why you don't hack into systems without permission. If you want to understand more how nuclear command, control, and decision making works, the two books to read are Command and Control: Nuclear Weapons, the Damascus Accident, and the Illusion of Safety by Eric Schlosser, and The 2020 Commission Report on the North Korean Nuclear Attacks Against the United States (A Speculative Novel) by Jeffrey Lewis.

Problem 7 Fail Caesar (12 points)

A student at a well known Junior University decided to write their own Caesar Cipher after learning about them in their computer security class. Unfortunately for the student, they fell asleep during the lecture on memory safety. (Note: The atoi() function converts the initial portion of the string to an integer, returning 0 in case of an error.)

```
1 #include < stdio.h>
  #include <stdlib.h>
   void encrypt(int offset, char plaintext[]) {
       char ciphertext [64];
 5
       memset (ciphertext, 0, 64);
6
       int i = 0;
 7
       fgets (plaintext, 64, stdin);
8
       while (plaintext[i]) {
9
           ciphertext[i] = plaintext[i] + offset;
10
11
12
       printf(ciphertext);
13
14
15
  int main(int argc, char *argv[]){
16
       char buffer [64];
17
       int offset = 0;
       if (argc > 1) offset = atoi(argv[1]) % 26;
18
19
       while (!feof(stdin)){
20
           memset (buffer, 0, 64);
21
           encrypt(offset , buffer);
22
23
       return 0;
24
```

- (a) What line contains a memory vulnerability? What is this vulnerability called?
- (b) Give a file that, when input to the command failcaesar with no arguments, will cause the program to crash.
- (c) How would you change the line to fix the vulnerability?
- (d) The student's friend who was awake for the memory safety lecture tells them to enable stack canaries to make their code more secure. If an attacker does not have time to perform a bruteforce attack, does enabling stack canaries prevent this code from being exploited? Explain why or why not.

#### Problem 8 A Lack of Integrity...

(9 points)

Alice and Bob want to communicate. They have preshared a symmetric key k. In order to send a message M to Bob, Alice encrypts it using AES-CBC, and sends the encryption to Bob. (You may assume that M's length is divisible by the AES-CBC block length and that characters are 8 bits, so no padding is necessary.) Recall that the actual message sent is IV||E(M), that is, the IV is prepended to the message and sent all as a single stream of bytes. Alice uses a random IV for each message.

In order to make sure that Bob is listening, they agree to using pingback messages. If Alice sends a message whose plaintext begins with the two bytes "PB", then Bob sends back the rest of the message  $in\ plaintext$ . For example, if Alice sends AES-CBC<sub>k</sub> ("PBI Love CS 161!"), then Bob responds "I Love CS 161!" without any encryption.

Alice uses the protocol to communicate some message M to Bob. Assume M is not a pingback message. Mallory, a man-in-the-middle attacker, decides to attempt to trick Bob into generating a pingback message. She thus sends the message IV'||IV||E(M), where IV' is a random 128b string.

- (a) With what probability will Mallory's message trigger a pingback message?
- (b) If Mallory's message triggers a pingback message, what does Mallory receive?
- (c) How can Alice and Bob change their protocol to prevent this attack?

Alice (who know send	n appropriate for the particular moves the plaintext of $M_1$ , that each block	using the same $\Gamma$ de) using AES (a ck of $M_1$ is different beknownst to Eve,	V/nonce and a deterministic padding scheme 128b block cipher). Eve, the Eavesdropper, at, that $M_1$ is 120 bytes, and that Alice never it turns out that the messages differ only in cal.
Yes,	Alice screwed up. But how badly? I	For each possibilit	y, select <i>all</i> which apply.
(a)	If Alice used AES-ECB (Electronic about $M_2$ :	Code Book), Eve	e is able to determine which of the following
	$\square$ That $M_2$ is exactly 120B long		That $M_2$ is less than 129B long but not the exact length
	$\square$ The entire plaintext for $M_2$		The plaintext for only the first two blocks of $M_2$
	$\square$ The entire plaintext for $M_2$ ex 2nd block	ccept for the	The plaintext for only the first block of $M_2$
(b)	If Alice used AES-CTR (Counter), 1	Eve is able to dete	ermine which of the following about $M_2$ :
	$\square$ That $M_2$ is exactly 120B long		That $M_2$ is less than 129B long but not the exact length
	$\square$ The entire plaintext for $M_2$		The plaintext for only the first two blocks of $M_2$
	$\square$ The entire plaintext for $M_2$ ex 2nd block	cept for the	The plaintext for only the first block of $M_2$
(c)	If Alice used AES-CBC (Cipher Bloabout $M_2$ :	ock Chaining), Even	e is able to determine which of the following
	$\square$ That $M_2$ is exactly 120B long		That $M_2$ is less than 129B long but not the exact length
	$\square$ The entire plaintext for $M_2$		The plaintext for only the first two blocks of $M_2$
	$\square$ The entire plaintext for $M_2$ ex 2nd block	ccept for the	The plaintext for only the first block of $M_2$
(d)	If Alice used AES-CFB (Ciphertex about $M_2$ :	t Feedback), Eve	is able to determine which of the following
	$\square$ That $M_2$ is exactly 120B long		That $M_2$ is less than 129B long but not the exact length
	$\square$ The entire plaintext for $M_2$		The plaintext for only the first two blocks of $M_2$
	$\square$ The entire plaintext for $M_2$ ex 2nd block	ccept for the	The plaintext for only the first block of $M_2$
(e)	If Alice did $not$ screw up, which model $M_3$ that is completely different from		etermine the exact length of a third message
	□ AES-ECB		AES-CBC

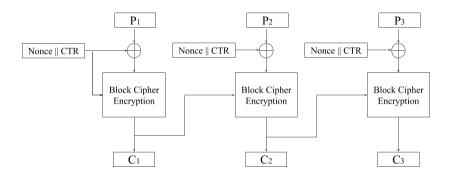
☐ AES-CFB

☐ AES-CTR

#### Problem 10 No More Keys

(7 points)

Frustrated by your newfound love of encryption schemes, your partner decides to throw away all of your secret keys. As a student in CS 161, you decide to make the best of a bad situation. You decide to design your own encryption scheme!



(a) Design the Decryption scheme.

(b) This is IND-CPA:

O True O False

(c) The encryption is parallelizeable:

O True O False

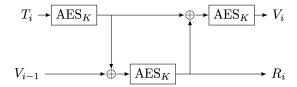
(d) The decryption is parallelizeable:

O True O False

#### Problem 11 Like Water off a DUHK's Back

(12 points)

The ANSI X9.17/X9.31 is a fairly simple pRNG that was widely used based on a block cipher (commonly AES). The internal state V and key K are combined with the current time T to update the state and produce a "random" value.



The current time is measured in microseconds as that is what the common operating system routines return. This is a strong pRNG as long as the initial state  $V_0$  and the key K are both high entropy and secret, and the block cipher is secure.

Unfortunately this scheme can fail badly when common mistakes are made. The standard never specified how to select K. So some implementations, rather than using a high-entropy source to seed a secret K, used a hardcoded key. The result is a catastrophic failure<sup>2</sup>.

- (a) If the attacker exactly knows K,  $T_1$ , and  $R_1$ , the attacker can then recover  $V_0$ . How?
- (b) Since one can then use this to calculate  $R_0$  given  $T_0$ , what design principle for a good pRNG does this fail to implement?
- (c) If the attacker knows  $T_0$  and  $T_1$  with just millisecond resolution, the attacker can check to see if a possible candidate for  $T_0$  and  $T_1$  is consistent with guesses for  $R_0$  and thereby know they found  $V_0$ . How many possible combinations of  $T_0$  and  $T_1$  may potentially need to be checked to determine  $V_0$ ?

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<sup>&</sup>lt;sup>2</sup>This was analyzed as the DUHK ("Don't Use Hardcoded Keys") attack, and it worked against FortiGate VPNs. For more details see https://duhkattack.com. This catastrophic failure mode is why it is no longer part of the standard suite of pRNGs.

# Foot-Shooting Prevention Agreement

I, \_\_\_\_ , promise that once

I see how simple AES really is, I will not implement it in production code even though it would be really fun.

This agreement shall be in effect until the undersigned creates a meaningful interpretive dance that compares and contrasts cache-based, timing, and other side channel attacks and their countermeasures.

