

# CS61C Sample Final

## 1 Opening a New Branch

```
struct market {
    char *name;
    struct item *inventory; // array of items
    int invSize; // size of array
};

struct item {
    int id;
    int price;
};

typedef struct market Market;
typedef struct item Item;
```

1) Consider the structures above which describe a Market and its inventory. Complete the function copy below, which makes a complete copy of the market structure and all of the data contained within it. None of the fields in the new copy should reference data in the original copy. Feel free to use the following functions:

```
char *strcpy(char *dst, const char *src);
size_t strlen(const char *str);
```

```
Market* copy(Market* orig) {
    Market* result =

    return result;
}
```

## 2 Thread Level Parallelism

For the following snippets of code below, Circle one of the following to indicate what issue, if any, the code will experience. Then provide a short justification. Assume the default number of threads is greater than 1. Assume no thread will complete before another thread starts executing.

Assume arr is an int array with length len.

a)

```
// Set all elements in arr to 0
int i;
#pragma omp parallel for
for (i = 0; i < len; i++)
    arr[i] = 0;
```

Sometimes incorrect

Always incorrect

Slower than serial

Faster than serial

b)

```
// Set element i of arr to i
#pragma omp parallel
for (int i = 0; i < len; i++)
    arr[i] = i;
```

Sometimes incorrect

Always incorrect

Slower than serial

Faster than serial

c)

```
// Set arr to be an array of Fibonacci numbers.
arr[0] = 0;
arr[1] = 1;
#pragma omp parallel for
for (int i = 2; i < len; i++)
    arr[i] = arr[i - 1] + arr[i - 2];
```

Sometimes incorrect

Always incorrect

Slower than serial

Faster than serial

### 3 Huge Pages

High performance applications commonly find themselves bogged down by the virtual memory subsystem when using large data sets. One solution is for the operating system to selectively allocate larger (huge) pages to back the virtual memory of such applications. We will look at one possible way these huge pages could be implemented. For this problem, we are using a 32-bit virtual and 32-bit physical address space, and want to support **8MiB huge pages**, with a normal page size of 4KiB.

To start off, we add a bit to each page table entry to indicate if it is a huge page. When this bit is set to 1, it means that the PPN is that of a 8MiB page in physical memory, and that this entry overrides the next 2047 entries for address translation. We require the virtual addresses of huge pages to be 8MiB aligned, so it's only legal for every 2048th page table entry to be marked as huge.

Example page table, where fields marked X are not used in address translation.

PPN	valid	dirty	huge
7342	1	0	0
6422	1	0	X (illegal)
1423	0	0	X (illegal)
...	...	...	...
9241	1	0	1
X (part of hugepage)	X	X	X (illegal)
X (part of hugepage)	X	X	X (illegal)

a) Fill out the number of bits needed for each field in the page table.

PPN:                  valid:                  dirty:                  huge:

b) With a fully-associative 16 entry TLB, what is the maximum amount of memory normally addressable with just cached page entries? What does this increase to if we allow huge pages? Write your answer in IEC format.

c) List a reason why huge pages might:

speed up translation	waste memory	slow down translation	be better for DMA

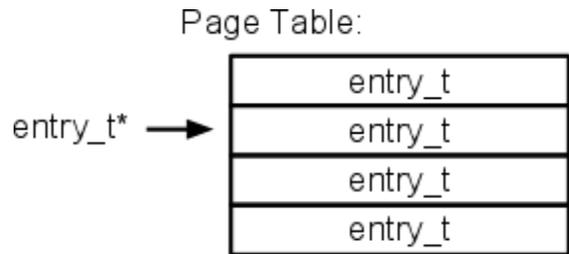
d) Describe what hardware modifications are required to allow for huge page support. In particular, what additional checks are necessary during a TLB lookup?

e) Since writing a huge page to disk can take a while, it is common for operating systems to split a huge page into smaller pages when swapping out such a page. Implement `split_hugepage()`, which is given a pointer to the huge page's entry. You can assume that the page table is an array of `entry_t`'s as defined below, and that the physical huge page can be divided up without moving it. The OS will take care of TLB invalidation after this function is called.

```
// definition of a page table entry
typedef struct entry {
    bool valid;
    bool huge;
    bool dirty;
    int ppn;
} entry_t;
```

```
void split_hugepage(entry_t *huge_ent) {
    huge_ent->huge = 0;
    // your code (HINT: huge_ent + i points to the ith entry after huge_ent)

}
```





d) You are trying to sell the value of this instruction. Give a situation in one sentence where this instruction would be valuable (would allow you to do the task faster than classical mips).

e) After you make these changes to the datapath, your boss tells you he now has to decrease the frequency of the clock for the CPU to work correctly. What did you do?

## 5 Synchronous Digital Circuits

a) You are an intern at a massive hardware firm. Your first task is to design a “prime number checker.” The circuit you must design will take as input an integer between 0 and 7. It will have three input bits, corresponding to the relevant integer. Your circuit should output 1 if the integer is prime and 0 otherwise. Neither zero nor 1 is prime.

Complete the Truth Table below. A, B and C correspond to the 4’s place, 2’s place and 1’s place respectively

Input A	Input B	Input C	Output X
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

b) Write the Boolean expression formed by the output in Sum-of-Products form. Reduce it to its simplest form.

c) Construct the circuit with the fewest gates using only AND, OR and NOT gates

d) Consider a system that can output a value between 0-3 with the ability to increment and decrement. This system will have two 1-bit inputs: increment and decrement and a 2-bit output (the count).

- If increment is high, the count should increase by one for the next cycle (wrap around if necessary).
- If decrement is high, the count should decrease by one for the next cycle (wrap around if necessary)
- If neither is high, the system should stay at the same value

- They will never both be high at the same time

Draw a state machine for this system with appropriate transitions for each input pair.

## 6 Who invited these people again?

Louis Reasoner writes the following self-modifying code:

```
foo:   la   $t0, modify
      sll  $a0, $a0, 11
      lw   $t1, 0($t0)
      addu $t2, $t1, $a0
      sw   $t2, 0($t0)
modify: addu $0, $0, $a1
      sw   $t1, 0($t0)
      jr   $ra
```

What happens if we call foo with \$a0=15 and \$a1=15? (3pt)

For each of the following questions, CIRCLE either a, b, c, d, or e.

Which set of inputs will permanently change the behavior of the program? (3pt)

- a) \$a0=3 and \$a1=3
- b) \$a0=5 and \$a1=5
- c) \$a0=7 and \$a1=7
- d) \$a0=9 and \$a1=9
- e) \$a0=11 and \$a1=11

Which set of inputs will always cause a bus error? (3pt)

- a) \$a0=7 and \$a1=7
- b) \$a0=8 and \$a1=7
- c) \$a0=7 and \$a1=8
- d) \$a0=8 and \$a1=8

Alyssa P. Hacker has figured out how foo works and wants to use the program for an unintended purpose – copying \$t3 to \$t4. Which set of inputs should she choose? (She could have just executed addu \$t4 \$0 \$t3 to achieve the same effect). (3pt)

- a) \$a0=0x000000CB and \$a1=0x000000CB
- b) \$a0=0x000000CC and \$a1=0x000000CC
- c) \$a0=0x000000CD and \$a1=0x000000CD
- d) \$a0=0x000000DB and \$a1=0x000000DB
- e) \$a0=0x000000DC and \$a1=0x000000D